

## PROOF FOUNDATIONS

### (W) WEISER DEFINITION OF SLICING:

Given a program  $P$ , an slicing criterion  $C=\langle v,s \rangle$  where  $v$  is a variable at statement  $s$ , and an slice  $S$ :  
If  $P$  halts on input  $I$ , then the value of  $v$  at statement  $s$  each time  $s$  is executed in  $P$  is the same in  $P$  and  $S$ . If  $P$  fails to terminate normally,  $s$  may be executed more times in  $S$  than in  $P$ , but  $P$  and  $S$  compute the same values for  $v$  each time  $s$  is executed by  $P$ .

### (A) DATA DEPENDENCE:

We say there exists a data dependence between two expressions when the first expression defines the value of a variable and the second one uses this value in at least one of the possible program executions without being any other expression modifying it.

NOTE: We consider that the arguments passed in a function call and the parameters of that function are a specific case of data dependence where the expression changes its name.

### (B) CONTROL DEPENDENCE:

There exists a control dependence between two expressions when the second expression cannot be evaluated without evaluating the first expression.

### (C) SEQUENTIAL REDUNDANCE:

When the return expression of a block or a function (the last expression of the block in Erlang) is a variable defined in the previous expression, this can be deleted avoiding the definition of this variable and returning the result of the previous expression, taking this expression the last position of the block and being returned in consecuese.

### (D) SYNTAX ERROR:

We say there exists a syntax error in a program when the removal or modification of a chosen expression transforms the program into a non-executable state.

### (E) SEMANTIC MODIFICATION:

There exists a semantic modification in an expression when the modification of one of its subexpressions modifies the behaviour of the whole expression.

### (F) ABSORBING PROPERTY:

A clause of a conditional or a function statement is absorbing when its guard is always evaluated to true or its pattern always matches.

### (G) FULL TEST VALIDATION:

There exists full test validation when an original program and a slice extracted from it can be executed with all possible input values of the original program and the values of the slicing criterion are the same in both executions.

NOTE: We consider in this definition also programs with slicing criteria that are independent of program inputs, where there is only one possible execution.

## COLOUR LEGEND

**Black:** Expressions deleted by executing phase 1 (iterative slicing with the selected slicers)

**Red:** Expressions deleted by executing phase 2 (modified ORBS algorithm)

**Green:** Expressions remaining in the quasi-minimal slices

**Orange:** Slicing Criterion

**NOTE1:** We will not prove whether black expressions of the program code can be deleted or not because they have been deleted by phase 1. Phase 1 produces a complete slice of the original code, so we can guarantee that these expressions are not part of the slice.

**NOTE2:** Our slices keep the syntax of the original program (we are not interested in amorphous slices). However, in order to make the final slice executable, some modifications of the source code are compulsory (e.g., replacing calls to deleted functions with a constant called "undef"). Therefore, we allow for some modifications of the source code to produce executable slices. The modifications made never affect the behaviour of the source code, they just ensure that the final code is a valid Erlang program.

```
%-----  
%-----  
%-- bench12.erl  
%--  
%-- AUTHORS:      Tamarit-Emartinm  
%-- DATE:         2013  
%-- PUBLISHED:    https://github.com/tamarit/edd/blob/master/examples/ternary (2016)  
%-- COPYRIGHT:    Bencher: The Program Slicing Benchmark Suite for Erlang  
%--               (Universitat Politècnica de València)  
%--               http://www.dsic.upv.es/~jsilva/slicing/bencher/  
%-- DESCRIPTION  
%-- The program performs translations and operations with balanced ternary. Its inputs are  
%-- two numbers in balanced ternary representation and a decimal number. It converts the  
%-- three numbers to the opposite representation and performs an operation with their  
%-- balanced ternary representation. The output is a tuple containing the input numbers  
%-- and the result in both representations.  
%-----  
%-----
```

```

-module(bench12).
-export([main/3]).

main(AS,B,CS) ->
  AT = from_string(AS), A = from_ternary(AT),
  BT = to_ternary(B),

  BS = to_string(BT),
  CT = from_string(CS), C = from_ternary(CT),
  RT = mul(AT,sub(BT,CT)),

  R = from_ternary(RT),
  RS = to_string(RT),
  [{AS,A},{BS,B},{CS,C},{RS,R}].

to_string(T) -> [to_char(X) || X <- T].

from_string(S) -> [from_char(X) || X <- S].

to_char(-1) -> $-;
to_char(0) -> $0;
to_char(1) -> $+.

from_char($-) -> -1;
from_char($0) -> 0;
from_char($+) -> 1.

to_ternary(N) when N > 0 ->

  to_ternary(N, []);

to_ternary(N) ->
  neg(to_ternary(-N)).

to_ternary(0,Acc) ->

  Acc;

to_ternary(N,Acc) when N rem 3 == 0 ->

```

%Given (A), B is necessary w.r.t.  $BT = to\_ternary(B)$   
%Given (A), BT is necessary w.r.t. the expression  $RT = mul(AT, sub(BT, CT))$   
%Replace  $to\_ternary(BT)$  with undef (NOTE2) would prevent to satisfy (1)&(2)  
%Delete  $RT = mul(AT, sub(BT, CT))$  would prevent to reach the SC because it is the only call to the sub function, where the SC is located  
%Replace  $sub(BT, CT)$  with undef would prevent to reach the SC  
%Given (A), BT is necessary w.r.t.  $sub(A, B)$ . Replace it with undef (NOTE2) would prevent to satisfy (1)&(2)

%Delete this clause would prevent to reach the SC  
%Given (A), N is necessary w.r.t.  $to\_ternary(N, [])$   
%Guard when  $N > 0$  cannot be deleted because it would prevent to satisfy (2)  
%Replace N with undef (NOTE2) would make this clause fulfill (F) and it would prevent to satisfy (2)  
%Replace 0 with undef (NOTE2) would prevent to satisfy (1) because of (E). This clause would become unreachable  
% $to\_ternary(N, [])$  cannot be deleted because it is the only expression of the clause. Replace it with undef (NOTE2) would prevent to satisfy (1)&(2)  
%Replace N with undef (NOTE2) would prevent to reach the SC because of a badarith error in the  $to\_ternary/2$  function  
%Replace [] with undef (NOTE2) would prevent to satisfy (1)&(2)  
%Delete this clause would prevent to satisfy (2)  
%Given (A), N is necessary w.r.t.  $neg(to\_ternary(-N))$   
% $neg(to\_ternary(-N))$  cannot be deleted because it is the only expression of the clause. Replace it with undef (NOTE2) would prevent to satisfy (2)  
%Replace  $to\_ternary(-N)$  with undef (NOTE2) would prevent to reach the SC because of a bad generator error in the neg function  
%Replace -N with undef (NOTE2) would prevent to reach the SC because of a badarith error in the  $to\_ternary/2$  function  
%This clause is the base case of the  $to\_ternary/2$  recursive function. Delete this clause would prevent to reach the SC because of an infinite loop  
%Replace 0 with \_ (NOTE2) would make this clause fullfils (F) and this would prevent to satisfy (1)&(2)  
%Given (A), Acc is necessary w.r.t. Acc  
%Acc is the only expression of the clause and it will be the returned value of the  $to\_ternary/2$  function because it is the returned value of its base case  
%Delete this clause would prevent to satisfy (1)&(2)  
%Given (A), N is necessary w.r.t.  $to\_ternary(N div 3, [0|Acc])$   
%Given (A), Acc is necessary w.r.t.  $to\_ternary(N div 3, [0|Acc])$   
%Delete guard when  $N rem 3 == 0$  would make this clause fullfils (F) and it would prevent to satisfy (1)&(2)  
%Replace  $N rem 3$  with undef (NOTE2) would prevent to satisfy (1)&(2) because this clause would be unreachable. This could be avoided by replacing also 0 with undef, but this would make this clause fullfils (F) and it would also prevent to satisfy (1)&(2)  
%Replace N or 3 in  $N rem 3$  with undef (NOTE2) would prevent to reach the SC because of a badarith error

```

to_ternary(N div 3, [0|Acc]);
to_ternary(N,Acc) when N rem 3 == 1 ->
to_ternary(N div 3, [1|Acc]);
to_ternary(N,Acc) ->
X = to_ternary((N+1) div 3, [-1|Acc]).

```

%Replace 0 with undef (NOTE2) would make this clause unreachable and it would prevent to satisfy (1)&(2). This could be avoided by replacing also N rem 3 with undef, but this would make this clause fullfils (F) and it would also prevent to satisfy (1)&(2)  
%to\_ternary(N div 3, [0|Acc]) cannot be deleted because it is the only expression of the clause. Replace it with undef (NOTE2) would prevent to satisfy (1)&(2)  
%Given (A), N div 3 is necessary w.r.t. to\_ternary(N,Acc). Replace N div 3 with undef (NOTE2) would prevent to reach the SC in (1)&(2) because of a badarith error  
%Replace N or 3 with undef (NOTE2) would prevent to reach the SC in (1)&(2) due to a badarith error  
%Given (A), [0|Acc] is necessary w.r.t. to\_ternary(0,Acc). This clause is the base case of the function and parameter Acc in the clause is necessary to calculate the returned value  
%Replace 0 with undef (NOTE2) would prevent to satisfy (1)&(2)  
%Delete Acc would prevent to satisfy (1)&(2) because each call to this clause would replace list Acc (returned in the base case) with list [0]  
%Delete this clause would prevent to satisfy (1)&(2)  
%Given (A), N is necessary w.r.t. to\_ternary(N div 3, [1|Acc])  
%Given (A), Acc is necessary w.r.t. to\_ternary(N div 3, [1|Acc])  
%Delete guard when N rem 3 == 1 would make this clause fullfils (F) and it would prevent to satisfy (1)&(2)  
%Replace N rem 3 with undef (NOTE2) would prevent to satisfy (1)&(2) because this clause would be unreachable. This could be avoided by replacing also 1 with undef, but this would make this clause fullfils (F) and it would also prevent to satisfy (1)&(2)  
%Replace N or 3 in N rem 3 with undef (NOTE2) would prevent to reach the SC because of a badarith error  
%Replace 1 with undef (NOTE2) would make this clause unreachable and it would prevent to satisfy (1)&(2). This could be avoided by replacing also N rem 3 with undef, but this would make this clause fullfils (F) and it would also prevent to satisfy (1)&(2)  
%to\_ternary(N div 3, [1|Acc]) cannot be deleted because it is the only expression of the clause. Replace it with undef (NOTE2) would prevent to satisfy (1)&(2)  
%Given (A), N div 3 is necessary w.r.t. to\_ternary(N,Acc). Replace N div 3 with undef (NOTE2) would prevent to reach the SC in (1)&(2) because of a badarith error  
%Replace N or 3 with undef (NOTE2) would prevent to reach the SC in (1)&(2) due to a badarith error  
%Given (A), [1|Acc] is necessary w.r.t. to\_ternary(0,Acc). This clause is the base case of the function and parameter Acc in the clause is necessary to calculate the returned value  
%Replace 1 with undef (NOTE2) would prevent to satisfy (1)&(2)  
%Delete Acc would prevent to satisfy (1)&(2) because each call to this clause would replace list Acc (returned in the base case) with list [1]  
%Delete this clause would prevent to satisfy (1)&(2) because of a matching error  
%Given (A), N is necessary w.r.t. to\_ternary((N+1) div 3, [-1|Acc])  
%Given (A), Acc is necessary w.r.t. to\_ternary((N+1) div 3, [-1|Acc])  
%X = to\_ternary((N+1) div 3, [-1|Acc]) cannot be deleted because it is the only expression of the clause. Replace this expression with undef (NOTE2) would prevent to satisfy (1)&(2)  
%Given (A), (N+1) div 3 is necessary w.r.t. to\_ternary(N,Acc). Replace (N+1) div 3 with undef (NOTE2) would prevent to reach the SC in (1)&(2) because of a badarith error  
%Replace (N+1) or 3 with undef (NOTE2) would prevent to reach the SC due to a badarith error  
%Replace N or 1 in (N+1) with undef (NOTE2) would prevent to reach the SC in (1)&(2) due to a badarith error  
%Given (A), [-1|Acc] is necessary w.r.t. to\_ternary(0,Acc). This clause is the base case of the function and parameter Acc in the clause is necessary to calculate the returned value  
%Replace -1 with undef (NOTE2) would prevent to satisfy (1)&(2)

%Delete Acc would prevent to satisfy (1)&(2) because each call to this clause would replace list Acc (returned in the base case) with list [-1]

```
from_ternary(T) -> from_ternary(T,0).
```

```
from_ternary([],Acc) ->
  Acc;
from_ternary([H|T],Acc) ->
  from_ternary(T,Acc*3 + H).
```

```
mul(A,B) -> mul(B,A,[]).
```

```
mul(_,[],Acc) ->
  Acc;
mul(B,[A|As],Acc) ->
  BP = case A of
    -1 -> neg(B);
    0 -> [0];
    1 -> B
  end,
  A1 = Acc++[0],
  A2=add(BP,A1),
  mul(B,As,A2).
```

```
neg(T) ->
  [-H || H <- T].
```

%Given (A), T is necessary w.r.t. [-H || H <- T]  
 %[-H || H <- T] cannot be deleted because it is the only expression of the function clause. Replace it with undef (NOTE2) would prevent to reach the SC in execution (2) due to a badarith error in the to\_ternary/2 function  
 %Replace -H with undef (NOTE2) would prevent to satisfy (2)

%Replace H with \_ (NOTE2) would prevent to reach the SC because of (D) w.r.t. [-H || \_ <- T]. This would be avoided by replacing -H with undef, but this would prevent to satisfy (2)

%Replace T with undef (NOTE2) would prevent to reach the SC due to a bad generator error in the list comprehension expression

```
sub(A,B) ->
  add(A,neg(B)).
```

%Given (A), A is necessary w.r.t. add(A,neg(B)) because it defines the value to the SC  
 %add(A,neg(B)) cannot be deleted because the SC would be deleted  
 %A cannot be deleted because it is the SC

```
add(A,B) when length(A) < length(B) ->
  add(lists:duplicate(length(B)-length(A),0)+A,B);
add(A,B) when length(A) > length(B) ->
  add(B,A);
add(A,B) ->
  add(lists:reverse(A),lists:reverse(B),0,[]).
```

```
add([],[],0,Acc) ->
  Acc;
add([],[],C,Acc) ->
  [C|Acc];
add([A|As],[B|Bs],C,Acc) ->
  [C1,D] = add_util(A+B+C),
  add(As,Bs,C1,[D|Acc]).
```

```
add_util(-3) -> [-1,0];
add_util(-2) -> [-1,1];
add_util(-1) -> [0,-1];
add_util(3) -> [1,0];
add_util(2) -> [1,-1];
add_util(1) -> [0,1];
add_util(0) -> [0,0].
```

#### EXECUTION RESULTS:

```
(1) B = 160
(2) B = -160
```

SLICING CRITERION  
 SC = [1,-1,0,0,-1,1]  
 SC = [-1,1,0,0,1,-1]